

A Wireless Sensor MAC Protocol for Bursty Data Traffic and Its Use in a Fire Monitoring Application

Authors:

Luis Bernardo - lfb@fct.unl.pt
Rodolfo Oliveira - rado@fct.unl.pt
Miguel Pereira - miguelpereira.pro@gmail.com
Ricardo Tiago - ricardo_tiago@hotmail.com
Miguel Silva - migueltsilva@gmail.com
Mário Macedo - mmm@fct.unl.pt
Paulo Pinto - pfp@fct.unl.pt

João Barros - barros@dcc.fc.up.pt

**SIGAPANO: Recolha de Informação Sensorial
com "Nós Patrulha"**

POSC/EIA/62199/2004



RTCM – Nov. 28, 2007

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

Outline

- Motivation and WSN MAC protocol brief overview
- MH-MAC presentation
- MH-MAC performance evaluation
- Alarm application overview
- Conclusions



Motivation

- Standard WSN MAC protocols are usually designed under the assumptions of **periodic traffic**, or **seldom traffic**, but not for applications where both characteristics are needed at different instants :
 - ◆ B-MAC and X-MAC were designed for seldom traffic:
 - ☞ Nodes run independent asynchronous duty cycles
 - ☞ Low Power Listening bind sender and receiver using a large preamble
 - ◆ S-MAC, T-MAC, and Z-MAC were designed for periodic traffic:
 - ☞ Nodes run synchronized duty cycles
 - ☞ T-MAC introduced abbreviated awake periods due to inactivity
 - ☞ Z-MAC changes to TDMA fallback during load peaks, maximizing throughput
 - ◆ **B-MAC/X-MAC maximum throughput is low**
 - ◆ **S-MAC/T-MAC/Z-MAC energy consumption for idle periods is high**

Motivation

- Application requirements change
 - ◆ e.g. Forest monitoring with mobile patrol nodes
 - ☞ Energy consumption must be minimized when the patrol node is away
 - ☞ Throughput must be maximized when the patrol node is present
 - ◆ e.g. Alarm application for large forest with mobile patrol nodes
 - ☞ The two requirements above ...
 - ☞ Delay must be minimized during a fire
- The existing protocols are too static
- ➔ **a new MAC protocol is needed to handle both periodic and seldom traffic, accordingly to the application needs**

Objectives

- Propose a **Multimode Hybrid MAC protocol (MH-MAC)** for packetizing radios (e.g. CC2420) that can be in one of three states:
 - ◆ **Asynchronous**
 - ◆ **Synchronous**
 - ◆ **Full-on**
- Provide applications an API to control the state
- Develop an alarm application using MH-MAC API
 - ◆ **Large forest** (e.g. Peneda-Gerês national park, with 72000 ha) **with sparse WSN coverage**

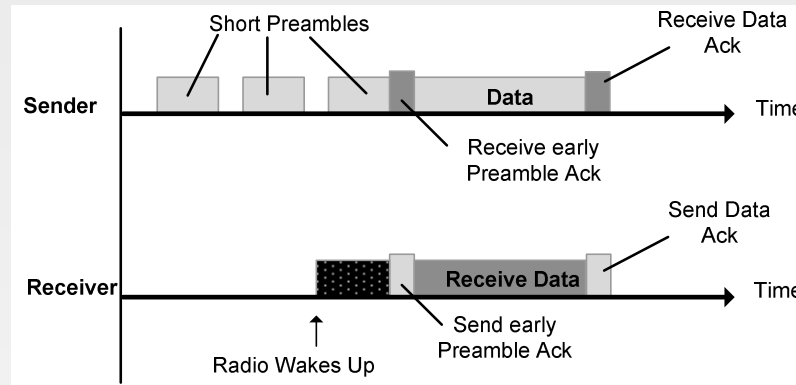
Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

MH-MAC asynchronous mode



Unicast asynchronous transmission

- Based on X-MAC unicast mode
- *Low Power Listening*
 - Sender sends a sequence of short preambles with a duration up to $2 * T_{DutyCycle}$ before the data frames
 - Receivers send an **Early Preamble ACK** when they awake up
- Collision receivers send **SHUT-UP packets**, with probability p

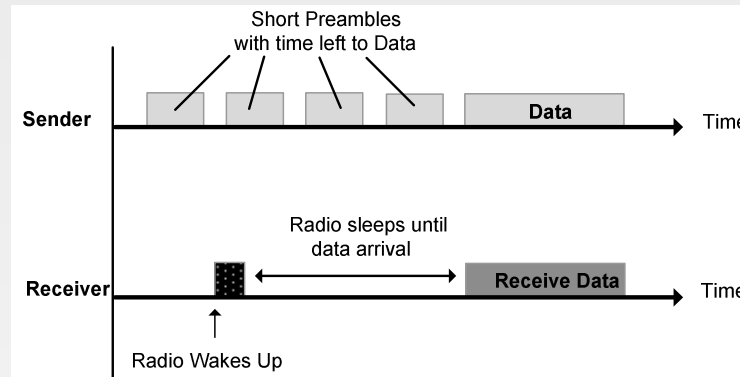
Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

MH-MAC asynchronous mode



Broadcast asynchronous transmission

- Sender sends a sequence of short preambles with a duration of $2x T_{DutyCycle}$ before the data frames
- **Receivers** receive the remaining preamble time in the preamble packet, and **sleep until near the data reception time**
- Collision receivers send SHUT-UP packets, with probability p

Outline

1. Motivation

2. MH-MAC Protocol

3. Performance Evaluation

4. Alarm Application

5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

MH-MAC synchronous mode

- Nodes run synchronized **duty-cycle periods**: a fixed length (11) sequence of slots.
- **Slots** have a fixed duration (100 ms)
 - Carry an average of 11,87 data frames (112 Bytes) on 802.15.4 radios
- Slots can be **public** or **private**
- **SYNC packets** synchronize the nodes and publicize the slot allocation map

Private slot

- **Reserved slot** for unicast traffic between two nodes.
- No collisions (differentiated backoffs)
- Unicast traffic is acknowledged
- SYNC packets include the local slot allocation map
- **After 25ms of inactivity** nodes go into sleep

Public slot (slot ALL)

- **Shared by all nodes**, used for broadcast traffic, or casual unicast traffic exchange.
- **First 15 ms reserved** for MAC signaling and application premium packets
 - SYNC packets
- **After 25ms of inactivity** nodes go into sleep
- Unicast traffic is acknowledged and run a contention-based protocol, preceded by a RTS/CTS exchange (above a threshold length).

Outline

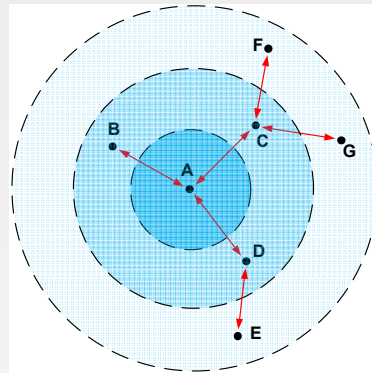
1. Motivation
2. **MH-MAC Protocol**
3. Performance Evaluation
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

MH-MAC synchronous mode

Example of a slot allocation



	Sinc	ALL	1	2	3	4	5	6	7	8	9	10
A												
B												
C												
D												
E												
F												
G												

- **Synchronous mode was designed for data collection over a sink tree**
- MH-MAC API only handles one hop interactions, but allows applications to create a multi-hop sink tree that
 - allocates multiple private slots to links proportional to the traffic
 - minimizes source-sink delay, organizing the private slots into a staggered wakeup schedule
- Delay is minimized in the figure
 - free slots are searched starting from the slot connecting to the sink
 - e.g. from F to A (slot 9 - 10)

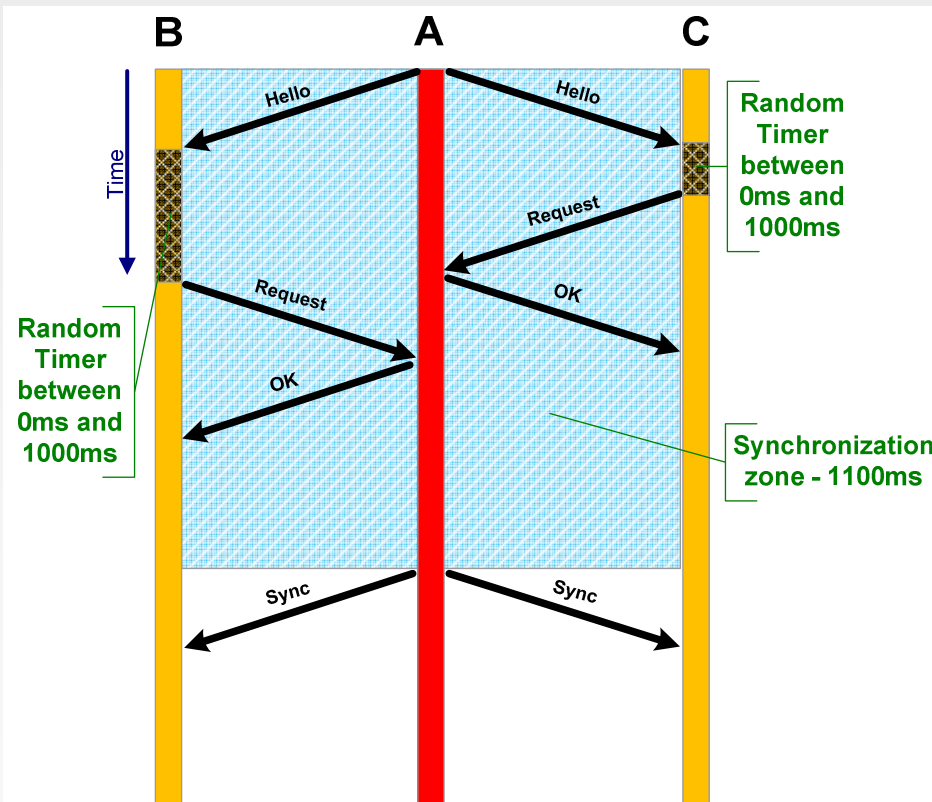
Outline

1. Motivation
2. **MH-MAC Protocol**
3. Performance Evaluation
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

MH-MAC mode swapping



Total synchronization time

- Tree shaped network
 - Depth r ; Neighbors $\leq v_{\max}$
- Time per node
 - Jitter before Hello $\leq T$
 - $2T$ preamble + 1100 ms
- Hypothesis:
 - Above 2 hops, nodes do not interfere
 - It may fail for normal topologies

$$TotalSyncTime \leq (2 + r(v_{\max} - 1)) \times (3T + 1100) [ms]$$

Outline

1. Motivation

2. MH-MAC
Protocol

3. **Performance
Evaluation**

4. Alarm
Application

5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

Performance evaluation

- We use TOSSIM simulator
 - ◆ Run MH-MAC nesC code
 - ◆ CC2420 radio stack modified
 - ◆ Additional meters measure active time/sleep time/tx time/rcv time
- Single hop scenarios:
 - ◆ multiple senders to a single receiver
 - ◆ 100 Bytes / packet
 - ◆ Variable average IPT (Inter Packet Time)
- Energy Estimation : Xbow Telos B current consumption

Operation	Current
Mote Standby (RTC on)	5,1 μ A
MCU Idle (DCO on)	54,5 μ A
MCU Active	1,8 mA
MCU + Radio RX	21.8 mA
MCU + Radio TX (0dBm)	19,5 mA

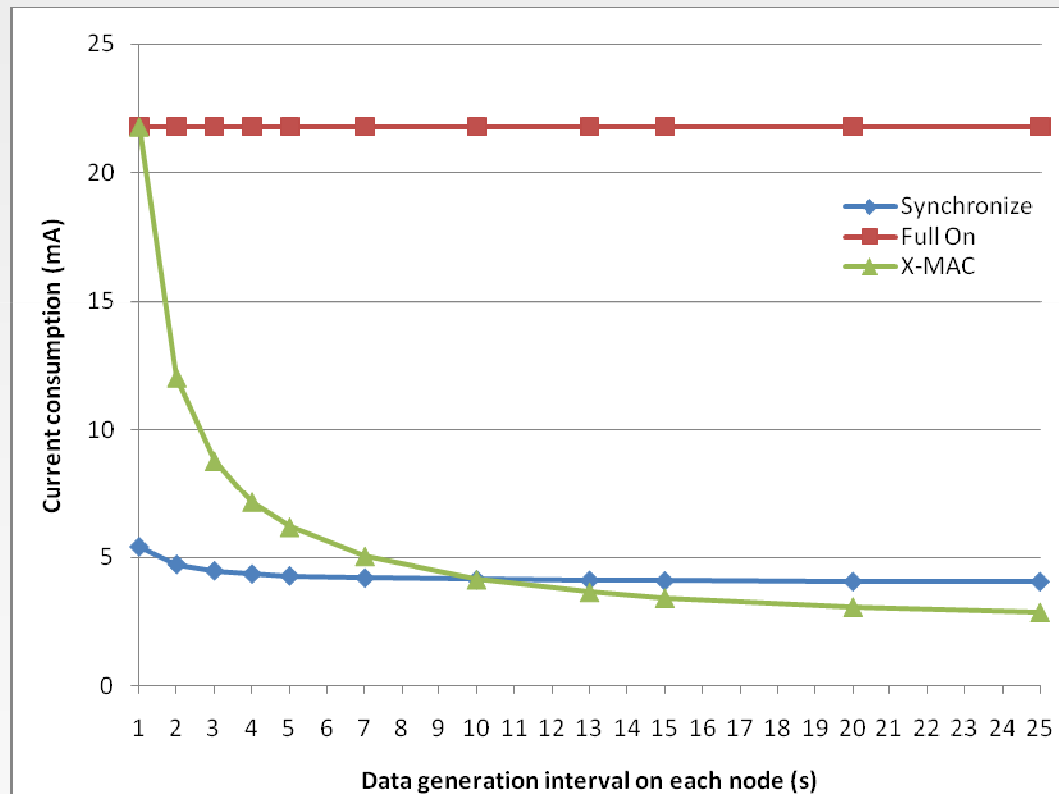
Outline

1. Motivation
2. MH-MAC Protocol
3. **Performance Evaluation**
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

Current consumption



- 4 senders / 1 receiver
- **Energy consumption is minimized for**
- *Asynchronous mode* (X-MAC) for IPT above 10 s
- *Synchronous mode* for IPT below 10s

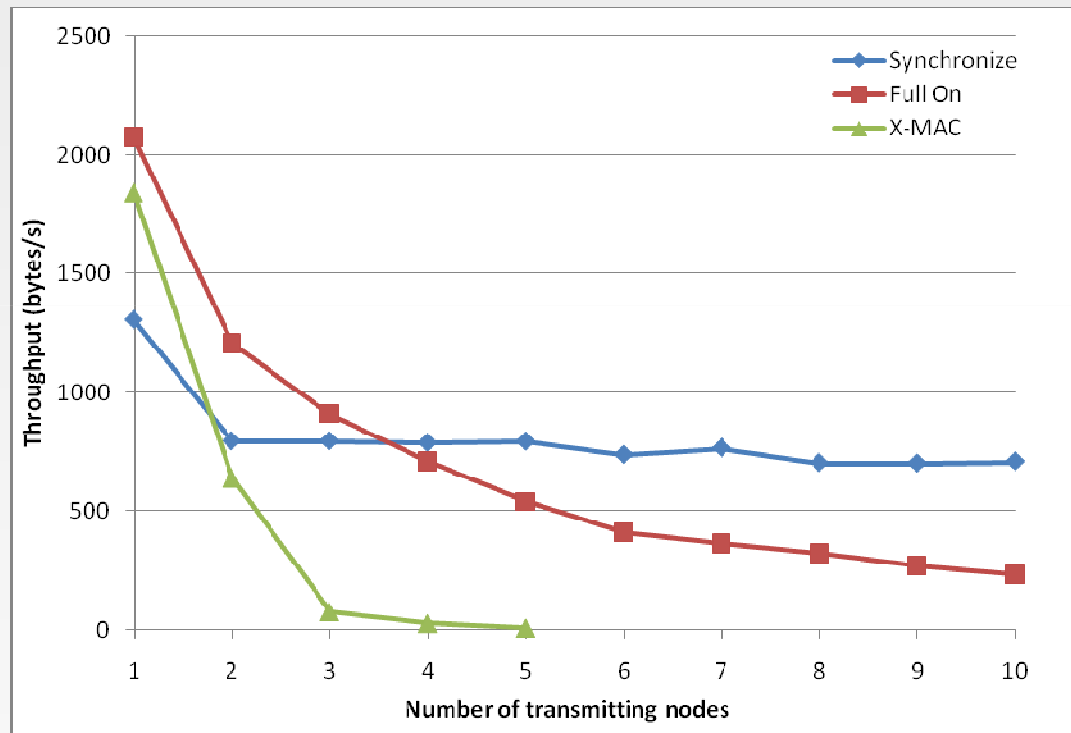
Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

Throughput



- N nodes transmit a burst of 20 packets, with IPT=0

- **Throughput is maximized for**

- *Full On mode* (no sleep) for $N \leq 3$

- *Synchronous mode* for $N > 3$

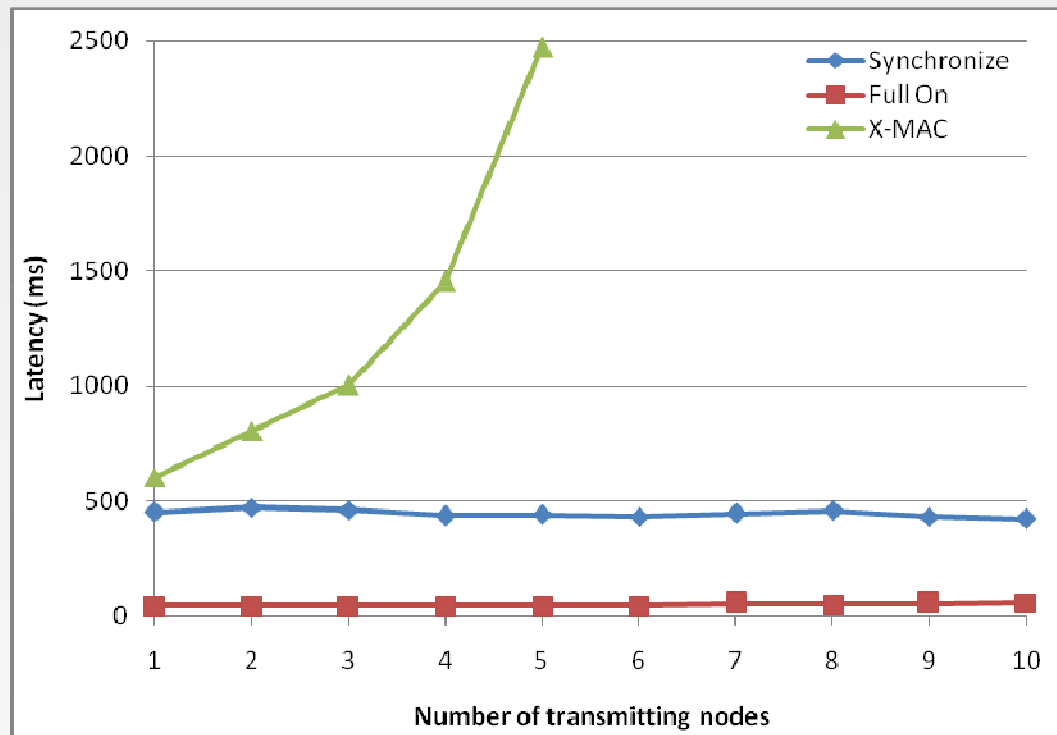
Outline

1. Motivation
2. MH-MAC Protocol
- 3. Performance Evaluation**
4. Alarm Application
5. Conclusions



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA

Latency



- $IPT=1s$
- Latency is minimized for
- *Full On mode* (no sleep)
- Asynchronous mode introduces a huge delay for more than 2 senders
- B-MAC/X-MAC reported good results only for 2 senders

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

Application overview

- Cross layering application, transport and MAC layer
 - ◆ Configure MH-MAC and transport level accordingly to the alarm level
- Mobile patrol nodes (laptop or PDA) roam the forest connecting to the scattered WSN:
 - ◆ Set and collect alarms
 - ◆ Motes run autonomously, when the patrol node is not present
- Two alarm thresholds
 - ◆ **Yellow alarm** – possible fire
 - ◆ **Red alarm** – imminent fire (high probability of mote destruction)
- Motes store the alarm record is safe motes (when possible)
 - ◆ Patrol nodes read the alarms from the storing motes

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

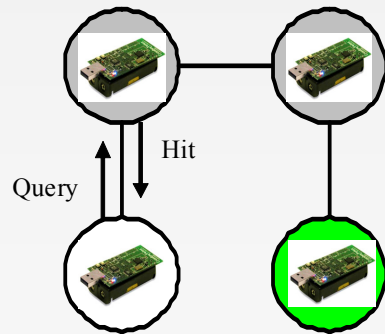
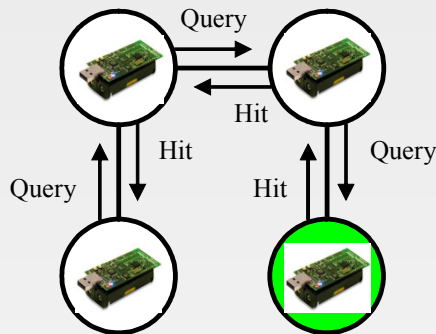
Application overview

- The application has three phases
 - ◆ Alarm deployment
 - ◆ Alarm handling
 - ◆ Alarm recovering
- Fast lookup operation for searching safe motes and stored records
 - ◆ Run a **peer-to-peer (p2p) lookup service** on the motes
 - **Clustering**
 - Run on a Virtual Overlay Network (VON) composed by a Minimum Connected Dominant Set (MCDS) covering the WSN
 - **QUERY flooding**
 - **HIT caching with invalidation mechanisms**

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

P2P searching algorithm

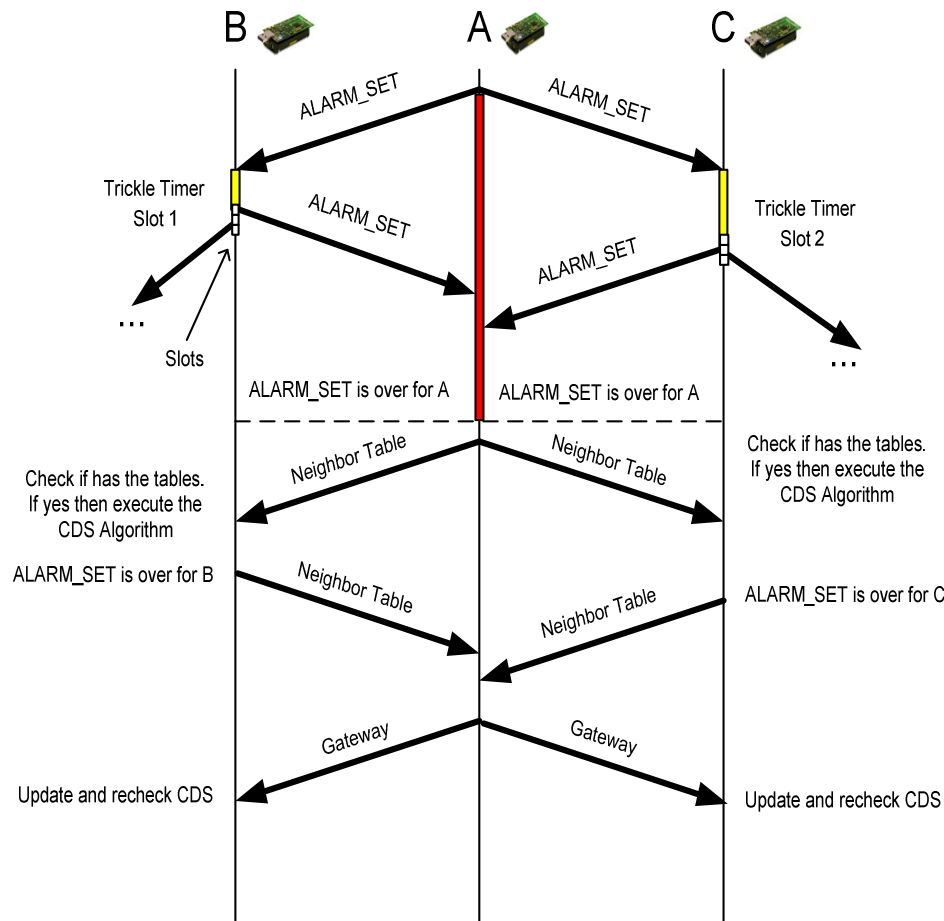


- Virtual Overlay Network of MCDS
 - MCDS motes act as ultra-peers
 - Classical Query flooding/Hit
 - Delay sensitive
 - Active Query Table
 - Cache of Hits forwarded
 - Forwarding motes p/Hit
- Hit carry two alarm levels
 - Sender alarm level
 - Maximum forwarding alarm level
- Forwarding nodes update the Hits when any of the alarm level changes

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. **Alarm Application**
5. Conclusions

Alarm deployment



Total deployment time

- Tree shaped network
 - Depth r ; Neighbors $\leq K$
- Time per node
 - T_{SLOT} per ALARM-SET
- Hypothesis:
 - No empty slots after trickle timer
 - Above 2 hops, nodes do not interfere

DeploymentTime $\geq$

$$(r-1)((2K-1)T_{SLOT} + T_{IDLE}) + 2T_{SLOT} + T_{IDLE}$$

Outline

1. Motivation

2. MH-MAC
Protocol

3. Performance
Evaluation

4. Alarm
Application

5. Conclusions

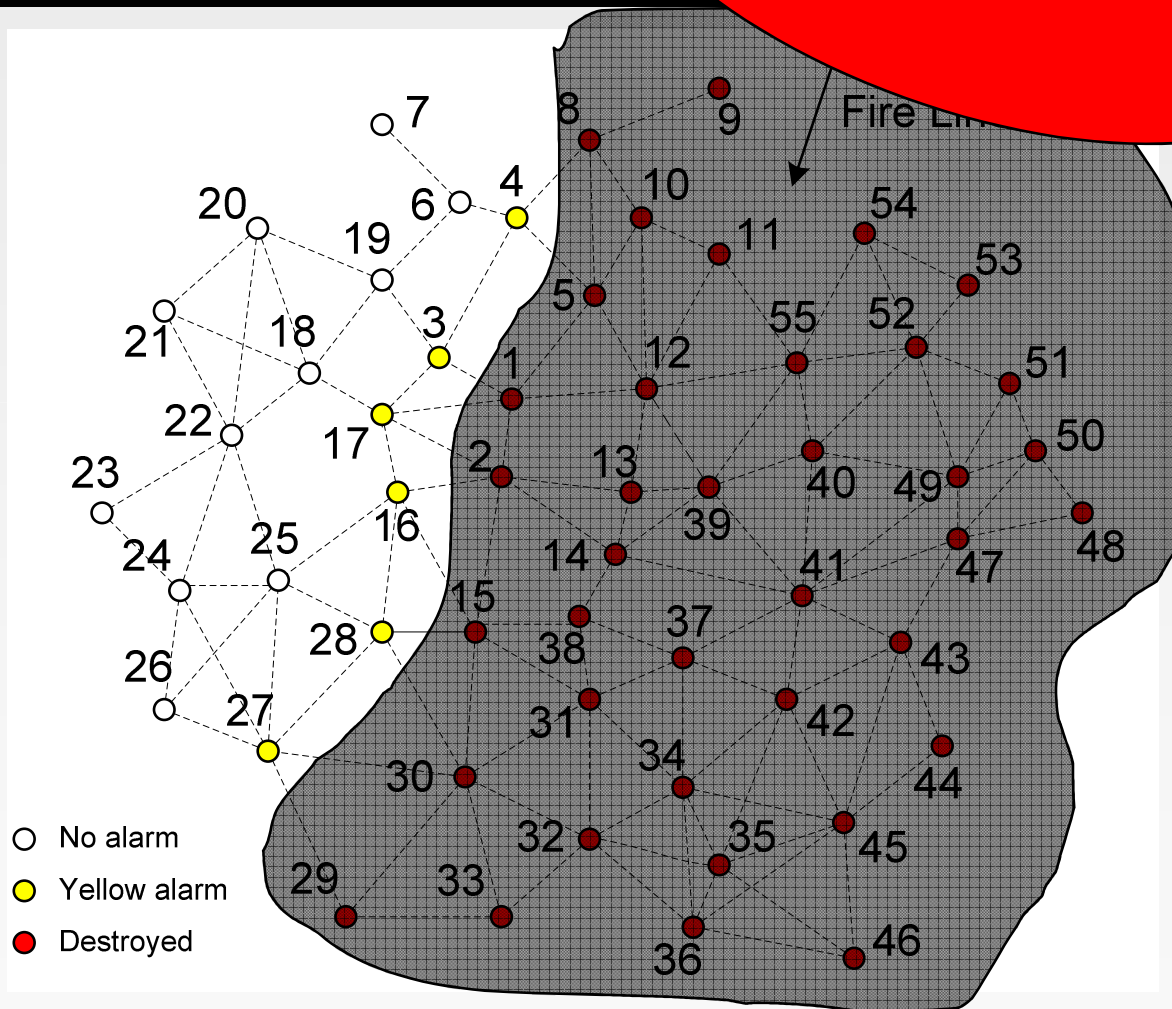
Alarm handling

- When an alarm is detected
 - MH-MAC changes to **Full-on**
 - Yellow alarm, red alarm and active CH motes
 - Motes record sensor measurement changes above a threshold limit to a repository mote located on a safe CH
 - return into no alarm state when the alarm thresholds are not exceeded for $T_{\text{NO-ALARM}}$ time
- Patrol nodes can receive ALARM-RECORD
 - asynchronously – CH store the ALARM-RECORD in the flash
 - real-time – CH forward the ALARM-RECORD it receives

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

Simulation scenario



- A fire front crosses a subset of the WSN

- Variable front velocity
- Constant temperature increase speed

- **Red alarm** (72 msgs)
 - $t_{yellow} + 120s$ – red
 - $t_{yellow} + 240s$ – destroy
- **Yellow alarm** (48 msgs)
 - during 360s

- **MEASUREMENT**
 - Percentage of saved ALARM-RECORD

Simulation Scenario

■ We evaluated three scenarios:

◆ **FullOn2Demand**

- ☞ MAC is asynchronous for no alarm situations and set to Full on when an alarm is detected

◆ **Synchronous**

- ☞ MAC is always in synchronous mode

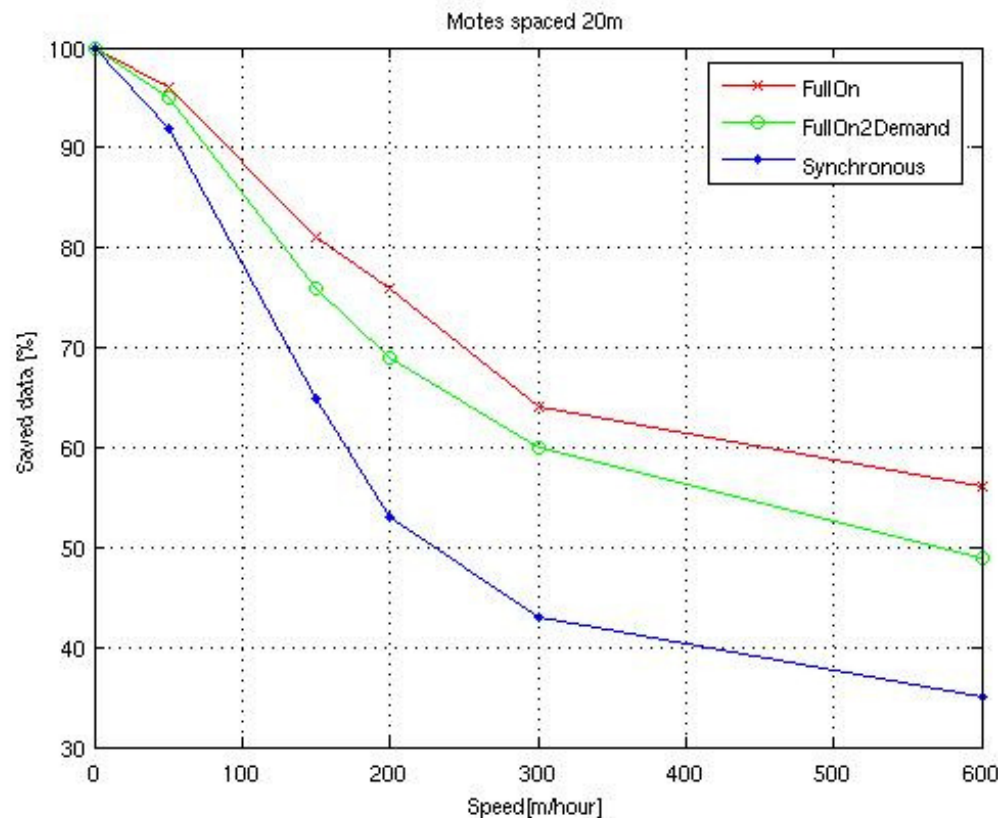
◆ **FullOn**

- ☞ MAC is always Full on (minimum delay)
 - a benchmark for evaluating the other two approaches

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

P2P without HIT invalidation

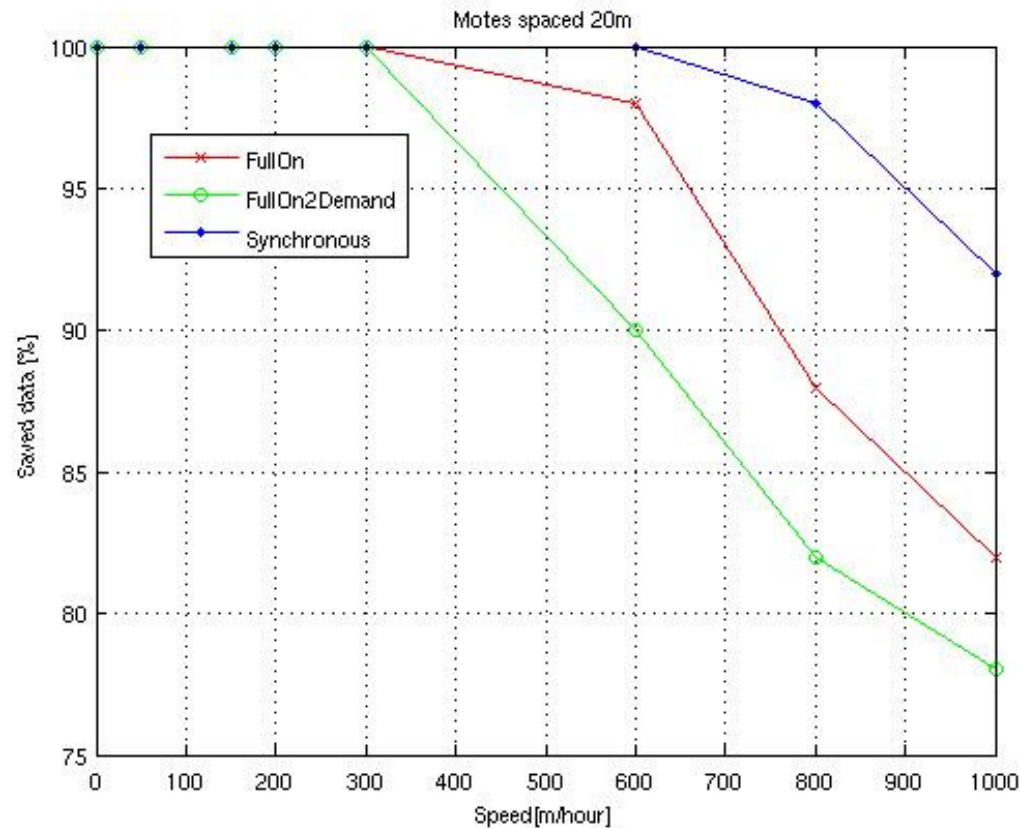


- loss occurs for 2.5 WSN hops / hour
- *Synchronous* fails because the p2p search delay is HIGH
 - delay per hop = $T_{DUTY_CYCLE}/2$
- *FullOn* fails because repository CH do not receive early warning of dangerous paths
- MH-MAC state transition is fast enough for the tested conditions
 - *FullOn2Demand* is slightly below *FullOn*

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Alarm Application
5. Conclusions

P2P with HIT invalidation



- P2P search is anticipated, recording paths are safer
 - loss occurs for 15 WSN hops / hour

- $I_{FullOn} = 21.80$ mA
- $I_{FullOn2Demand} = 10.62$ mA
- $I_{Sync} = 10.78$ mA
- SYNC costs on idle time

- Synchronous maximizes throughput (**no collisions**)
- MH-MAC state transition is fast enough for the tested conditions
 - *FullOn2Demand* is slightly below *FullOn*

Conclusions

- MH-MAC allows applications to choose the mode, with three possible objectives:
 - **Asynchronous mode** for minimizing energy consumption;
 - **Synchronous mode** for maximizing throughput for convergecast traffic
 - **Full On mode** for minimizing latency
- State transition incurs in energy costs and delay – applications must use it conscientiously
 - Alarm application shows that it can be used in real-time scenarios

Future Work

- WSN mobility support (Mobile MH-MAC is being developed)
 - Improves radio interference handling and reduced state transition overhead
- Thorough testing of MH-MAC and the applications on a physical WSN

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Conclusions

More information can be downloaded from:

<http://tele1.dee.fct.unl.pt>

L. Bernardo, R. Oliveira, M. Pereira, M. Macedo, P. Pinto, “A Wireless Sensor MAC Protocol for Bursty Data Traffic”, in IEEE PIMRC'07, September 2007.

L. Bernardo, R. Oliveira, R. Tiago, P. Pinto, “A Fire Monitoring Application for Scattered Wireless Sensor Networks: A peer-to-peer cross-layering approach”, in WINSYS'07 (part of ICETE'07), July 2007.

Outline

1. Motivation
2. MH-MAC Protocol
3. Performance Evaluation
4. Conclusions

Thanks for your attention.

Q&A ...



FACULDADE DE
CIÊNCIAS E TECNOLOGIA
UNIVERSIDADE NOVA DE LISBOA